

An efficient non-condensed approach for model predictive control

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Optimizing Cyber Physical Systems

Outline

- Optimization problem formulation
- **Equality constraint elimination** and proposed **nonlinear least-squares** approach with **bounded variables**
- Bounded-variable least-squares (**BVLS**) solver
- Problem **sparsity** and **matrix abstraction** (**build-free MPC**)
- Numerically-stable **sparse recursive QR factorization**
- Numerical results

Discrete-time nonlinear model

General multivariable nonlinear (NL) prediction model

$$\mathcal{M}(Y_k, U_k, S_k) = \mathbf{0}$$

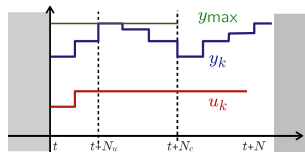
- Past inputs $U_k = (u_{k-n_b}, \dots, u_{k-1})$, $u_k \in \mathbb{R}^{n_u}$
- Past outputs $Y_k = (y_{k-n_a}, \dots, y_k)$, $y_k \in \mathbb{R}^{n_y}$
- Measured exogenous signals $S_k = (s_{k-n_c}, \dots, s_{k-1})$, $s_k \in \mathbb{R}^{n_s}$
- n_a, n_b and n_c define the model order
- Special case (state-space model): $U_k = u_k$, $Y_k = x_k$
- Assumption: \mathcal{M} is differentiable
- Examples: NL state-space models, deterministic **parameter-varying** NL-ARX models (black-box), I/O difference equations from first principles, **neural networks** with smooth activation function...
- On linearization about arbitrary \hat{U}, \hat{Y} :

$$- A(S_k)_0 \Delta y_k = \sum_{j=1}^{n_a} A(S_k)_j \Delta y_{k-j} + \sum_{j=1}^{n_b} B(S_k)_j \Delta u_{k-j} + \mathcal{M}(\hat{Y}, \hat{U}, S_k),$$

A, B represent required Jacobians

MPC problem setup

- Prediction horizon N , control horizon N_u
- $z_k = \{u_k, \dots, u_{k+N_u-1}, y_{k+1}, \dots, y_{k+N}\}$
=vector of decision variables



- Performance index

$$\min_z \|f_k(z)\|_2^2$$

Examples:

- f_k is linear, $f_k(z) = W_k(z - z_{\text{ref},k})$ (standard tracking problem)
- f_k is arbitrary nonlinear differentiable function
- Constraints
 - (nonlinear) equality constraints due to the prediction model \mathcal{M}
 - **upper** and **lower bounds** on **inputs** and **outputs** $p_k \leq z \leq q_k$
 - general **inequality constraints** $g(u_{k+j}, y_{k+j}) \leq 0$ can be **softened** and treated as equalities $g(u_{k+j}, y_{k+j}) + \sigma_j = 0$, with $\sigma_j \geq 0$

Constrained nonlinear programming (NLP) problem

- Consider tracking problem with quadratic costs for simplicity (everything immediately extends to arbitrary nonlinear costs $\|f_k(z)\|_2^2$)

Resulting NLP formulation at each sample step k

$$\begin{aligned} \min_z \quad & \frac{1}{2} \|W_k(z - z_{\text{ref},k})\|_2^2 \\ \text{s.t.} \quad & h_k(z, \phi_k) = \mathbf{0}, \\ & p_k \leq z \leq q_k. \end{aligned}$$

- Matrix W is often diagonal and the Jacobian of $h(z)$ is sparse and structured
- Initial condition vector ϕ consists of past I/O values

Proposed NMPC formulation

Key Idea

Soften equality constraints via quadratic penalties

$$\min_{p_k \leq z \leq q_k} \frac{1}{2} \|W_k(z - z_{\text{ref},k})\|_2^2 + \frac{\rho}{2} \|h_k(z)\|_2^2$$

- Penalty parameter $\rho > 0$ is a large weight
- **Motivation:** model is uncertain anyway, so why impose $h_k(z, \phi_k) = 0$ exactly?
- The problem can be rewritten as

$$\min_{p_k \leq z \leq q_k} \frac{1}{2} \|r_k(z)\|_2^2, \quad r_k(z) = \begin{bmatrix} \frac{1}{\sqrt{\rho}} W_k(z - z_{\text{ref},k}) \\ h_k(z, \phi_k) \end{bmatrix}$$

- Box-constrained nonlinear least squares problem is **always feasible**
- **Fast solution** using bounded-variable nonlinear least squares (**BVNLLS**)
- **Same control performance** as with conventional NMPC/NLP

Bounded-variable nonlinear least squares (BVNLLS)

Problem: **Sum-of-squares** cost function with **box constraints**

$$\min_{p_k \leq z \leq q_k} \frac{1}{2} \|r_k(z)\|_2^2$$

- **Gauss-Newton method**: efficiently solves **unconstrained** nonlinear LS
 - Sequence of linear least-squares problems
 - Hessian (H) is approximated as $H \approx J^\top J$, $J = \nabla_z r(z)^\top$
 - Rapid convergence with good initial guess
 - Only **first-order information** (J) is needed
- Proposed solver: Gauss-Newton method with box constraints
- Line-search problem: **Linear least-squares** with box constraints (BVLS)
- **Guaranteed convergence** using sufficient decrease condition

Bounded-variable nonlinear least squares (BVNLLS)

Problem: Sum-of-squares cost function with box constraints

$$\min_{p_k \leq z \leq q_k} \frac{1}{2} \|r_k(z)\|_2^2$$

- BVNLLS: Gauss-Newton method with box constraints

Bounded-variable nonlinear least squares

Initialize $z^{(0)} \in \{z | p \leq z \leq q\}$, $j \leftarrow 0$

- 1: Update Jacobian $J \leftarrow \nabla_z r(z^{(j)})^\top$; (**Linearization**)
- 2: Compute gradient $J^\top r$ of Lagrangian function;
- 3: If first-order optimality conditions are satisfied then **stop**;
- 4: Solve $\Delta \leftarrow \arg \min_{p - z^{(j)} \leq \Delta \leq q - z^{(j)}} \|J\Delta + r(z^{(j)})\|_2^2$ via **BVLS solver**; (Line search)
- 5: Compute step-size $0 < \alpha \leq 1$ for backtracking;
- 6: $z^{(j+1)} \leftarrow z^{(j)} + \alpha\Delta$; (Update iterate)
- 7: $j \leftarrow j + 1$; **go to** Step 1;

- **Linear MPC** case exactly recovered by **single BVNLLS iteration**!

Bounded-variable least squares (BVLS)

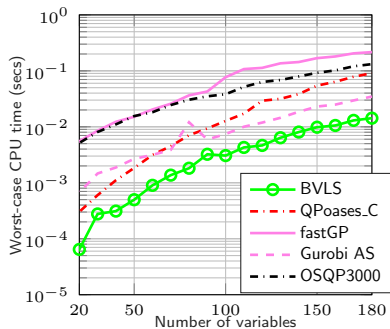
Problem: Least-squares with box constraints

$$\min_{p \leq x \leq q} \frac{1}{2} \|Jx - b\|_2^2$$

- BVLS [2] is a primal **active-set algorithm**
- Finds solution x^* by iterating until the optimal active set (\mathcal{A}^*) is found
- Active set? $\mathcal{A} = \{i | x(i) = p(i)\} \cup \{i | x(i) = q(i)\}$
- Main computations:
 - Solve unconstrained LS: $J(:, \neg \mathcal{A})^\dagger (b - J(:, \mathcal{A})x(\mathcal{A}))$ (every iteration)
 - Gradient entries: $J(:, \mathcal{A})^\top (b - Jx)$ (in some iterations)
- \mathcal{A} is updated by one index (inserted or removed)
 - \Rightarrow Subsequent LS problems are related by insertion or deletion of **1** column in J !

BVLS solver

- BVLS solves a sequence of related LS problems
- Efficient implementation with numerically stable **recursive QR updates** [3]
- Library free, simple arithmetic operations
- Stable also in **single precision**
- Suitable for **embedded hardware platforms**
- Implemented and tested on a **real industrial PLC** (paper under preparation)



Double precision, random **poorly-conditioned**, box-constrained LS problems (2.6GHz Intel Core i5 Mac)

BVLS for MPC: Problem Sparsity

Problem: Least-squares with box constraints

$$\min_{p \leq x \leq q} \frac{1}{2} \|\mathbf{J}x - b\|_2^2$$

- $J = \begin{bmatrix} W_k \\ \nabla_z h_k(z_k, \phi_k)^\top \end{bmatrix}$
- Problem can be **constructed** using sequence of affine models obtained from linearization over previously computed or guess trajectory

$$-A(S_k)_0^{(i)} \Delta y_k = \sum_{j=1}^{n_a} A(S_k)_j^{(i)} \Delta y_{k-j} + \sum_{j=1}^{n_b} B(S_k)_j^{(i)} \Delta u_{k-j} + \mathcal{M}(\hat{Y}^{(i)}, \hat{U}^{(i)}, S_k)$$

(i = prediction step)

$$x = \Delta z_k = \{\Delta u_k, \Delta y_{k+1}, \dots, \Delta u_{k+N_u-1}, \Delta y_{k+N_u}, \Delta y_{k+N_u+1}, \dots, \Delta y_{k+N}\}$$

- Outputs are kept as decision variables (non-condensed approach) for a larger but sparse problem formulation which is cheap to construct

BVLS for MPC: Problem Sparsity

Problem: Least-squares with box constraints

$$\min_{p \leq x \leq q} \frac{1}{2} \|\mathbf{J}x - b\|_2^2$$

$$\nabla_z h_k(z_k, \phi_k)^\top = \left[\begin{array}{cccccc|cccccc} B_1^{(1)} & A_0^{(1)} & \mathbf{0} & \mathbf{0} & \dots & & & & & & \mathbf{0} \\ B_2^{(2)} & A_1^{(2)} & B_1^{(2)} & A_0^{(2)} & \mathbf{0} & \dots & & & & & \mathbf{0} \\ & \vdots & & \ddots & & & & & & & \vdots \\ B_{N_u}^{(N_u)} & A_{N_u-1}^{(N_u)} & & & & \dots & B_1^{(N_u)} & A_0^{(N_u)} & \mathbf{0} & \dots & \mathbf{0} \\ \hline B_{N_u+1}^{(N_u+1)} & A_{N_u+1}^{(N_u+1)} & & \dots & B_3^{(N_u+1)} & A_2^{(N_u+1)} & \sum_{i=1}^2 B_i^{(N_u+1)} & A_1^{(N_u+1)} & A_0^{(N_u+1)} & \mathbf{0} & \dots & \mathbf{0} \\ B_{N_u+2}^{(N_u+2)} & A_{N_u+1}^{(N_u+2)} & \ddots & & B_4^{(N_u+2)} & A_3^{(N_u+2)} & \sum_{i=1}^3 B_i^{(N_u+2)} & A_2^{(N_u+2)} & A_1^{(N_u+2)} & A_0^{(N_u+2)} & \mathbf{0} & \dots & \mathbf{0} \\ & \vdots & & \ddots & \ddots & & \vdots & & \vdots & & \ddots & & \mathbf{0} \\ B_{N_p}^{(N_p)} & A_{N_p-1}^{(N_p)} & & \dots & B_{N_p-N_u+2}^{(N_p)} & A_{N_p-N_u+1}^{(N_p)} & \sum_{i=1}^{N_p-N_u+1} B_i^{(N_p)} & A_{N_p-N_u}^{(N_p)} & A_{N_p-N_u-1}^{(N_p)} & \dots & A_1^{(N_p)} & A_0^{(N_p)} \end{array} \right]$$

$$z = \{u_k, y_{k+1}, \dots, u_{k+N_u-1}, y_{k+N_u}, y_{k+N_u+1}, \dots, y_{k+N}\}$$

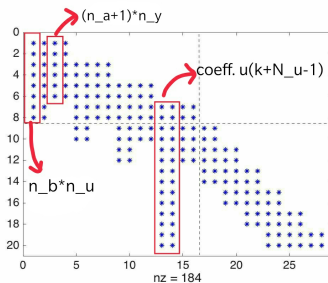
Structure depends on the ordering of decision variables, model (n_a, n_b, n_u, n_y) and tuning parameters (N, N_u)

BVLS for MPC: Problem Sparsity

Problem: Least-squares with box constraints

$$\min_{p \leq x \leq q} \frac{1}{2} \|Jx - b\|_2^2$$

Structure depends on the ordering of decision variables, model and tuning parameters



Sparsity pattern of $\nabla_z h_k(z)$ for a random model with $N_p = 10$, $N_u = 4$, $n_a = 2$, $n_b = 4$, $n_u = 2$ and $n_y = 2$.

Build-free MPC

- Typical MPC setup:
 - Step 1: Construct an optimization problem based on the prediction model and tuning parameters (e.g., N, N_u)
 - Step 2: Pass it in *standard* form to an optimization solver
- Constructing optimization problem matrices can be time consuming, especially in approaches with *condensed* formulation
- A change in the model coefficients, horizons, tuning weights, model size, needs re-construction of the optimization problem
- We propose methods that systematically eliminate the problem construction phase, resulting in:
 - reduced memory requirement
 - faster execution
 - ability to adapt to changes in model/tuning parameters at runtime at no computational cost

Key Idea

Parameterize the optimization algorithm in terms of model and tuning parameters

Matrix abstraction

BVLS problem: Least-squares with box constraints

$$\min_{p \leq x \leq q} \frac{1}{2} \|\mathbf{J}x - b\|_2^2$$

- The sparse Jacobian matrix $J = \begin{bmatrix} W_k \\ \nabla_z h_k(z_k, \phi_k)^\top \end{bmatrix}$ contains tuning weights and coefficients from the sequence of affine models with **indexing completely defined by model and tuning parameters**
- Role of J in BVLS:
 - Solve unconstrained LS: $J(:, \neg \mathcal{A})^\dagger (b - J(:, \mathcal{A})x(\mathcal{A}))$
 - Gradient entries: $J(:, \mathcal{A})^\top (b - Jx)$
- **Key observation:** all operations with J can be replaced by 2 abstract operators
 - 1) **Jix**: return i th column of J times a given scalar x
 - 2) **JtiX**: return i th column of J times a given vector X (i th column of $J = i$ th row of J^\top)

Two operators to replace all J instances

- To code **Jix** and **JtiX** we need
 - 1) Model coefficients from the sequence of affine models: store all in a single vector **M** (faster execution) **or** generate online via linearization routines (lower memory)
 - 2) Model parameters n_a, n_b, n_u, n_y
 - 3) Tuning parameters N, N_u and weights
- For J times a vector v , use **Jix** over each element of v and accumulate the result
- For J^\top times a vector v , use **JtiX** over each element of v and store result in the corresponding element of output vector
- Recall: location of non-zeros is already known in terms of model and tuning parameters!
 - \implies Only non-zero entries in J are operated
 - \implies Matrix operations as fast as sparse linear algebra while using significantly lesser amount of memory!

Build-free MPC algorithm

BVNLLS without problem construction

Initialize $z^{(0)} \in \{z | p \leq z \leq q\}$, $j \leftarrow 0$

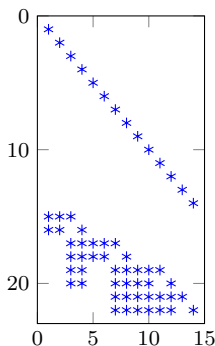
- 1: ~~Update Jacobian~~ $J \leftarrow \nabla_z r(z^{(j)})^\top$ **Update M**; (**Linearization**)
- 2: Compute gradient $J^\top r$ of Lagrangian function; **Use JtiX**
- 3: If first-order optimality conditions are satisfied then **stop**;
- 4: Solve $\Delta \leftarrow \arg \min_{p - z^{(j)} \leq \Delta \leq q - z^{(j)}} \|J\Delta + r(z^{(j)})\|_2^2$ via **BVLS solver**; **Uses Jix, JtiX**
- 5: Compute step-size $0 < \alpha \leq 1$ for backtracking;
- 6: $z^{(j+1)} \leftarrow z^{(j)} + \alpha\Delta$; (Update iterate)
- 7: $j \leftarrow j + 1$; **go to** Step 1;

Code of **Jix** and **JtiX** does not change with any change in model or tuning parameters or problem size

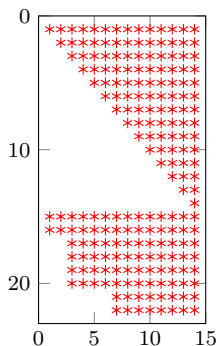
\Rightarrow entire MPC code is stand-alone for a given problem formulation

Sparse matrix factors

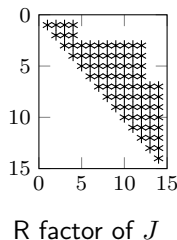
- Recall: we solve $J(:, \neg \mathcal{A})^\dagger (b - J(:, \mathcal{A})x(\mathcal{A}))$ in each BVLS iteration
- Best way: **recursive thin QR factorization** (using Gram-Schmidt orthogonalization)
- How to exploit sparsity? How to know the location of non-zeros in QR?



Sparse matrix $J = QR$



Q factor of J



R factor of J

Sparsity pattern of Jacobian J and its thin QR factors for a random NARX model with diagonal weights, and parameters $n_y = n_u = 2$, $n_a = 2$, $n_b = 1$, $N = 4$, $N_u = 3$.

Gram-Schmidt orthogonalization

- If $J = QR$,

$$Q'(:, i) = J(:, i) - \sum_{j=1}^{i-1} Q(:, j) Q(:, j)^{\top} J(:, i),$$

$$Q(:, i) = Q'(:, i) / \left\| Q'(:, i) \right\|_2.$$

$$R(j, i) = Q(:, j)^{\top} J(:, i), \forall j \in [1, i-1],$$

$$R(i, i) = \left\| Q'(:, i) \right\|_2$$

- Above formulae refer to **classical** Gram-Schmidt process: catastrophic numerical cancellation possible
- We use the theoretically equivalent **modified** Gram-Schmidt method with **automatic reorthogonalization** for numerical stability

Sparsity analysis

- Define the non-zero structure of a vector x to be the set of indices $\mathcal{S}(x)$ such that $x(i) \neq 0, \forall i \in \mathcal{S}(x)$, and $x(j) = 0, \forall j \notin \mathcal{S}(x)$.

Theorem: Non-zero structure of columns of Q factor

Consider an arbitrary sparse matrix $J \in \mathbb{R}^{n_1 \times n_2}$ of full rank such that $n_1 \geq n_2$ and let Q denote the Q-factor from its thin QR factorization i.e., $J = QR$. The non-zero structure of each column $Q(:, i)$ of Q satisfies

$$\mathcal{S}(Q(:, i)) \subseteq \bigcup_{j=1}^i \mathcal{S}(J(:, j)), \forall i \in [1, n_2],$$

$$\text{and } \mathcal{S}(Q(:, 1)) = \mathcal{S}(J(:, 1)).$$

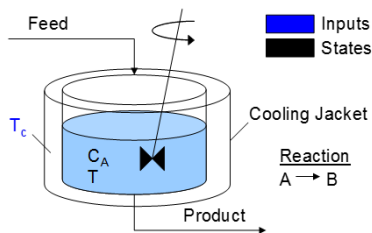
- Using the above theorem, which is based on MGS, and its corollaries [4], we exploit sparsity without even storing J !

(paper [4] = Saraf, Bemporad 2019 available on arXiv)

Recursive QR updates

- With diagonal weights, sparsity pattern info of Q factor can be stored in just 2 integer vectors
- Recursive update in sparsity pattern \implies update entries in 2 vectors of dimension = no. of columns of J
- Main principle: thin QR factorization of a matrix is **unique**
 \implies column indices of J may be added to or removed from the active set in an arbitrary order!
- R factor's sparsity exploited using orthogonality: $R = Q^\top J$

Numerical example: NMPC of CSTR



Continuous Stirred Tank Reactor (CSTR) ¹

● CSTR model

$$T^{(k+1)} = T^{(k)} + t_s(T_f^{(k)} - 1.3T^{(k)} + \kappa_1 C_A^{(k)} e^{\frac{-5963.6}{T^{(k)}}} + 0.3T_c^{(k)})$$

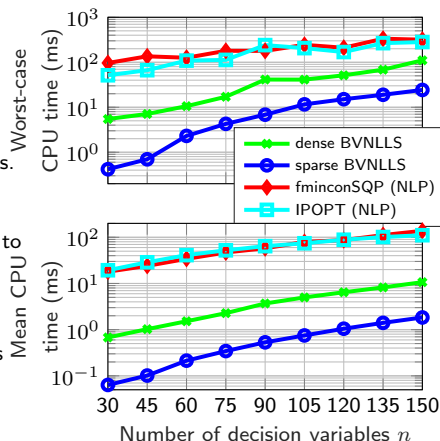
$$C_A^{(k+1)} = C_A^{(k)} + t_s(C_{Af}^{(k)} - \kappa_2 C_A^{(k)} e^{\frac{-5963.6}{T^{(k)}}} - C_A^{(k)})$$

- Nonlinear system with 2 outputs, 1 input and 2 measured disturbances
- Model coefficients as in MPC Toolbox demo (Mathworks)

¹ retrieved from apmonitor.com

Execution time: Small-sized problems

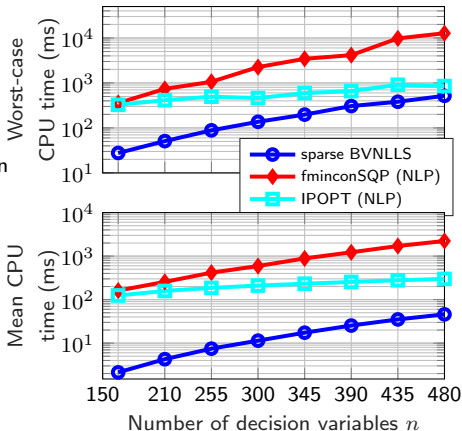
- **SQP**: subproblems may be infeasible.
Warm start exploited (MATLAB `fmincon`)
- **IPOPT**: uses MA57 solver, sparse routines.
No warm start exploited
- **BVNLLS**: very few Gauss-Newton steps to converge, exploits warmstarts
- Sparse linear algebra (LA) is typically slower than dense LA for small problems
 - proposed sparsity exploiting methods allow $\approx 10\times$ faster solution than dense variant even for small problem sizes!
 - $\approx 100\times$ speedup on mean CPU time w.r.t. benchmarks



Solver comparison in MATLAB for NMPC of CSTR simulated on a Mac with 2.6GHz Intel Core i5.
 $N = n/3$, no. of box constraints = n pairs, no. of equality constraints = $2N$, $\sqrt{\rho} = 10^4$,
 $N_{\text{sim}} = 1500$ sample steps.

Execution time: Larger problems

- Active-set methods can be faster than interior-point methods if sparsity is exploited [1]
- Faster even for large problems
- sparse BVNLLS tool for NMPC
 - Efficient C implementation
 - easily embeddable



Solver comparison in MATLAB for NMPC of CSTR simulated on a Mac with 2.6GHz Intel Core i5. $N = n/3$, no. of box constraints = n pairs, no. of equality constraints = $2N$, $\sqrt{\rho} = 10^4$, $N_{\text{sim}} = 1500$ sample steps.

Conclusions

- **Key ideas:**
 - **relax equality constraints** due to dynamics using penalty functions
 - **parameterize optimization solver** in terms of MPC parameters
- The proposed optimization solvers for MPC are:
 - **simple** to code, **fast** to execute, **flexible** in real time
 - good for **embedded platforms** (PLCs, μ controllers, ...)
 - **competitive** with state-of-the-art algorithms
- **Novel** linear algebra methods devised to heavily exploit **sparsity**
- Unifying MPC framework for LTI/LPV/NLTI/NLPV systems
- Linearization step can be code-generated using symbolic math software, no other code generation required - **easy deployability!**
- Extensions: Matrix-free MPC, more general problems

Key References



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